

RSA program Report / Started running on: Thu Nov 26 21:58:55 2009

TRACE OF GENERATING SAMPLE INTEGER

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random number : 0 is: 25024
In binary it is: 110000111000000
The least significant bit of that number is obviously ==> 0
So therefore bit number 1 from left is: 0

random number : 1 is: 28864
In binary it is: 111000011000000
The least significant bit of that number is obviously ==> 0
So therefore bit number 2 from left is: 0

random number : 2 is: 25294
In binary it is: 110001011001110
The least significant bit of that number is obviously ==> 0
So therefore bit number 3 from left is: 0

random number : 3 is: 1535
In binary it is: 1011111111
The least significant bit of that number is obviously ==> 1
So therefore bit number 4 from left is: 1

random number : 4 is: 4316
In binary it is: 1000011011100
The least significant bit of that number is obviously ==> 0
So therefore bit number 5 from left is: 0

32-bit Padded random integer: 00000000000000000000000001000101

THIS IS A SAMPLE FOR AN UNSUCCESSFUL PRIMALITY CHECK
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Checking the integer: 115 , Using a= 58

i	X[i]	z	Y	Y
6	1	1	1	58
5	1	58	29	72
4	1	72	9	62
3	0	62	49	49
2	0	49	101	101
1	1	101	81	98
0	0	98	59	59

 The integer: 115 is not a prime!

THIS IS A SAMPLE FOR A SUCCESSFUL PRIMALITY CHECK (PERHAPS PRIME)
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Checking the integer: 101 , Using a= 79

i	X[i]	z	Y	Y
6	1	1	1	79
5	1	79	80	58
4	0	58	31	31
3	0	31	52	52
2	1	52	78	1
1	0	1	1	1
0	0	1	1	1

 The integer: 101 is perhaps a prime!

Our n=8989 Our Phi(8989)=8800

Checking if e=5 can be a public key [Extended Euclidean Algorithm]

q	r1	r2	r	s1	s2	s	t1	t2	t
1760	8800	5	0	1	0	1	0	1	-1760

Unfortunately e=5 Does not have a multiplicative inverse in $Z(8800)$, so we go ahead and try the next e

Checking if e=6 can be a public key [Extended Euclidean Algorithm]

q	r1	r2	r	s1	s2	s	t1	t2	t
1466	8800	6	4	1	0	1	0	1	-1466
1	6	4	2	0	1	-1	1	-1466	1467
2	4	2	0	1	-1	3	-1466	1467	-4400

Unfortunately e=6 Does not have a multiplicative inverse in $Z(8800)$, so we go ahead and try the next e

Checking if e=7 can be a public key [Extended Euclidean Algorithm]

q	r1	r2	r	s1	s2	s	t1	t2	t
1257	8800	7	1	1	0	1	0	1	-1257
7	7	1	0	0	1	-7	1	-1257	8800

FOUND IT ! e will be:7 Its multiplicative inverse in $Z(8800)$ is:7543
So, ==> d=7543

HERE IS Alice'S CRYPTO SYSTEM

As Integers:

Alice's 'n' in 32 bits is:000000000000000000010001100011101

K ==> Bob found first '1' from left of 'n' in position ==> K= 13

So the u composed by Bob (according to the specs):

u[b]=00000000000000000001011110011111

This u as integer is =====> u =6047

So Bob sends to Alice u= 6047

And then Alice [after hashing and decrypting with her private key] returns to Bob v= 603

Showing Bob's encryption of V E(V, e)

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i	X[i]	Squaring Y	Multiplying Y
3	1	$1^2 \pmod{8989} = 1$	$603 \times 1 \pmod{8989} = 603$
2	1	$603^2 \pmod{8989} = 4049$	$603 \times 4049 \pmod{8989} = 5528$
1	1	$5528^2 \pmod{8989} = 5173$	$603 \times 5173 \pmod{8989} = 136$

Bob encrypts v with Alice's public key and gets Z= 136

Bob also does hash of: 6047 which in binary is:

00000000000000000001011110011111 and he gets h(u)= 136

AND IF THEY MATCH [Z == h(u)], BOB KNOWS HE IS TALKING TO SOMEONE WHO'S GOT ALICE'S PRIVATE KEY (HOPEFULLY ALICE)

TO SUMMARIZE

In Integers:

u = 6047
h(u) = 136
V=D(d,h(u)) = 603
E(e,V) = 136

In Binary:

u = 000000000000000000001011110011111
h(u) = 00000000000000000000000010001000
V=D(d,h(u)) = 0000000000000000000001001011011
E(e,V) = 00000000000000000000000010001000